

The Drop Shunting Problem

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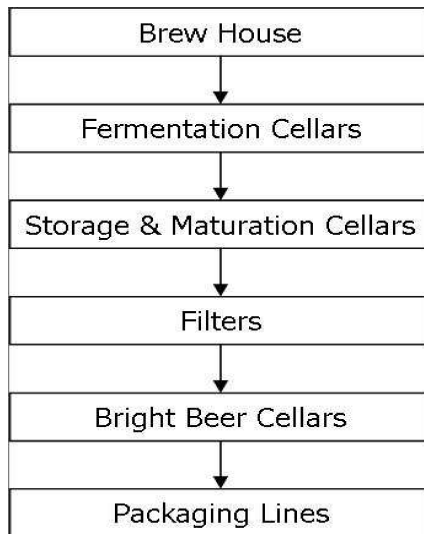
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Outline

- 1 Maco View of Brewerys
- 2 The Conceptual Layout
- 3 Demand, Stock Keeping Units (SKU), Packaging Types
- 4 Efficient Delivery by Trailer-Truck for Meeting Daily Demand
- 5 Heuristic Method of Solution
- 6 The Mathematical Modelling Approach

- The brewhouse (BH),
- Fermentation (FV),
- Storage and maturation (SV),
- Filtration,
- Bright beer cellars (BBT),
- Packaging lines.

Macro View of a Brewery



The Conceptual Layout

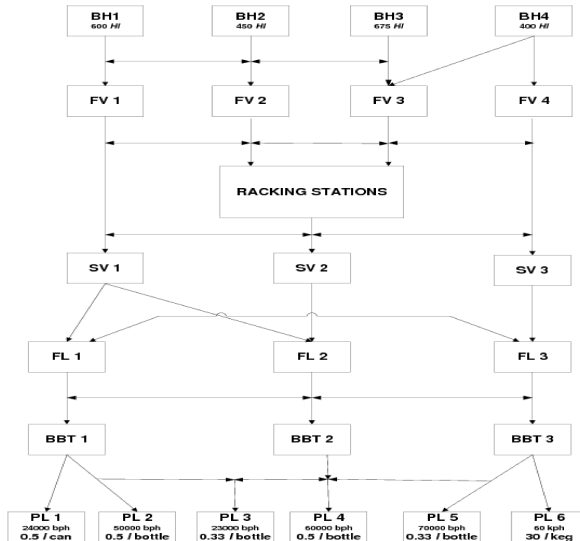


Table: Examples of SKUs

Brand	Pack Type	Demand at t_o
<i>A</i>	0.33 / Bottle	3897 <i>HI</i>
<i>A</i>	0.5 / Can	3122 <i>HI</i>
<i>B</i>	0.33 / Bottle	4865 <i>HI</i>
<i>B</i>	0.5 / Can	3028 <i>HI</i>
<i>E</i>	0.33 / Bottle	1465 <i>HI</i>
<i>E</i>	0.5 / Bottle	4144 <i>HI</i>
<i>E</i>	0.5 / Can	1012 <i>HI</i>

Table: Pack Types

Pack Type	Description	PTV_j (l)
1	Keg	30
2	Bottle	0.33
3	Can	0.5
4	Bottle	0.5

Brand $A = 1$, Pack Type=2 \Rightarrow SKU=12

A Weekly Scheduling Problem

#	SKU	Brand	Pack Type	Volume (Hl)
1	11	1	1	180
2	14	1	4	12782
3	23	2	3	3602
4	24	2	4	13574
5	32	3	2	10959
6	44	4	4	865
7	54	5	4	3732
8	64	6	4	2371
9	71	7	1	370
10	73	7	3	3674
11	74	7	4	6802
12	81	8	1	580
13	83	8	3	3616
14	84	8	4	7415
15	92	9	2	1506

The Conceptual Layout

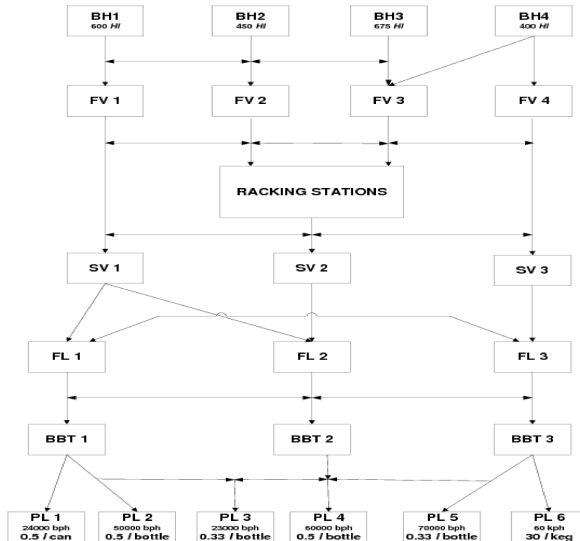


Figure: Demand Met By Efficient Use of Trucks and Trailers



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Assumptions

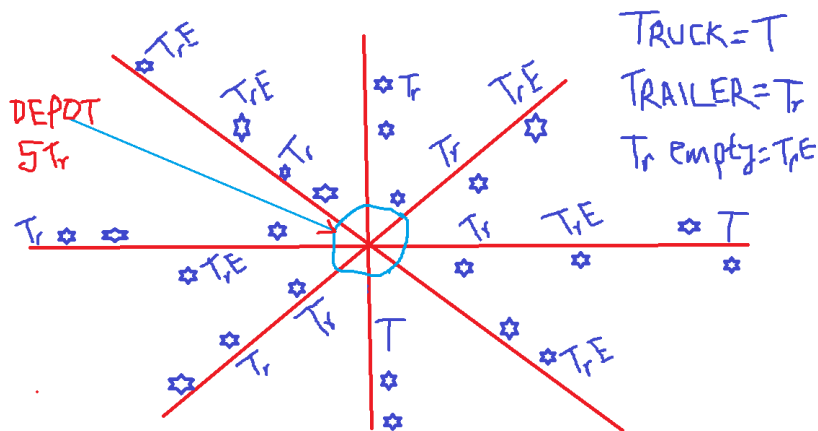
- A single Depot is considered. A geographical Area Filled by Customers
- $\#Trucks=p$; $\#Trailers=q$; $\#Customers=r$; $p \ll q \ll r$.
- Demands are uploaded in Trailers and Trucks just moves Trailers around (or drop them to the Depot).
- All Trailers are uniform (all have the same Capacity).
- Capacity of the Trailer: Multiples of average demand per customer.
- Demands are such that a Trailer can Upload the maximum demand.
- A Trailer moves back and forth (from Customer to Depot)–Multiple Uploading
- Activities are Daily

- Customer 1
 - SKU=21: 25,000 Bottle
 - SKU=13: 20,000 Cans
 - ...
- Customer 2
 - SKU=34: 15,000 Bottle
 - SKU=33: 10,000 Cans
 - SKU=22: 5,000 Bottles
 - SKU=13: 5,000 Cans

Table: Pack Types

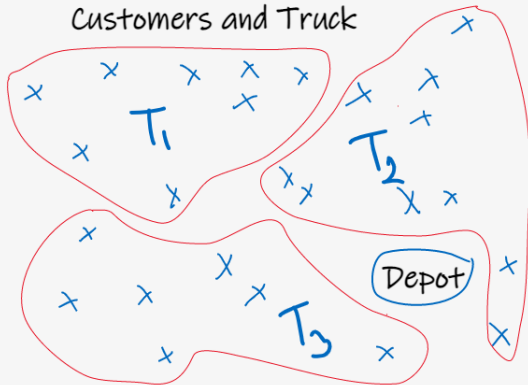
Pack Type	Description	$PTV_j (l)$
1	Keg	30
2	Bottle	0.33
3	Can	0.5
4	Bottle	0.5

Figure: Location of Demands, Trucks and Trailers



*=Customer; T =Truck, Tr =Trailer, TrE =Empty Trailer

Distributions of Customers and Truck



- Each Trailer: job_i, job_j, job_k (Trip₁); job_k, job_l, job_m (Trip₂) and so on
- Which Sequence is better for the Truck:

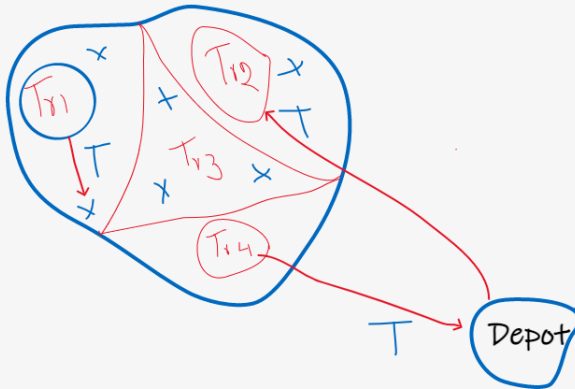
Truck : $(Tr_1, \cdot), (Tr_2, \cdot), (Tr_3, \cdot), (Tr_4, \cdot), (Tr_5, \cdot), (Tr_1, \cdot), (Tr_2, \cdot),$
 $(\cdot, TrE_1), (Tr_6, \cdot) \dots$

Truck : $(Tr_1, \cdot), (Tr_2, \cdot), (Tr_1, \cdot), (Tr_2, \cdot), (Tr_1, \cdot), (Tr_2, \cdot), (\cdot, TrE_1),$
 $(\cdot, TrE_1), (Tr_3, \cdot) \dots$

Distance Travelled by a Truck is dependent of the job-I assigned to Trailers in a cluster.

T: Truck; Tr: Trailer;
x: Customer

Cluster of Customers



- Subdivide the Customers into p number of Clusters (there are p Trucks): Use the p -mean algorithm
- Each Truck is assigned to one cluster; Assign a number of Trailers at each Cluster (e.g. q/p Trailers per Cluster)
- Consider Each Demand as a Job (Job-I); Job-II are assignment of $T_r E$ (empty trailers) and non empty trailers to a Truck.
- Upload a number of Type-I jobs in a Cluster to each Trailer for each Trip. A sequence of Uploads is needed for Each Trailer (**Optimal Decision is needed here**).
- Each Truck needs to perform a sequence of Type-II jobs (**Optimal Sequencing is needed here**)
- Daily Demand Scenario (Created Randomly) **Must Be Met**

- $(\# \text{Trucks}, \# \text{Trailers}, \# \text{Customers}) = (p, q, r)$; Depot: $DP \in \mathbb{R}^2$;
- Customer Locations: $\{C(1), C(2), \dots, C(r)\}$, $C(i) \in \mathbb{R}^2$;
- Demands: $\{D(1), D(2), \dots, D(r)\}$, $D(i) > 0$
- An Optimal Ordered Partition $\{R^1, R^2, \dots, R^p\}$ s.t.
 $\{C(1), C(2), \dots, C(r)\} = \bigcup_i^p R^i$, $R^i \cap R^j = \emptyset$
- The Ordered Set $R^k = \{C(r_1^k), C(r_2^k), C(r_3^k), \dots, C(r_s^k)\}$
- Cost Function for Truck k :

$$c(R^k) = d(DP, C(r_1^k)) + d(C(r_1^k), DP) + d(DP, C(r_2^k)) \\ + \sum_{i=1}^{s-1} d(C(r_i^k), C(r_{i+1}^k)) + d(C(r_s^k), DP) + \dots$$

$$(Tr_1, \cdot), (Tr_2, \cdot), (Tr_1, \cdot), (Tr_2, \cdot), (\cdot, TrE_1), \dots$$

$$x_{ij}^k = \begin{cases} 1 & \text{if } i\text{-th job-I}(i\text{-th Demand) is assigned to Trailer } j \text{ in Cluster } k \\ 0 & \text{otherwise} \end{cases}$$

$$\min \sum_k^p c(R^k) \quad s.t. \quad (1)$$

$$\sum_i^r \sum_j^q \sum_k^p x_{ij}^k = \sum_i^r D(i) \quad (2)$$

where

$$c(R^k) = d(DP, C(r_1^k)) + d(C(r_1^k), DP) + d(DP, C(r_2^k)) \\ + \sum_{i=1}^{s-1} d(C(r_i^k), C(r_{i+1}^k)) + d(C(r_s^k), DP) + \dots$$

Thank You
